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Core stable algorithms for coalition games with complementarities and peer effects

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\section{ABSTRACT}
In this short paper, we show two new algorithms for finding stable structures in ordinal coalition potential games. The first one is anytime and enumerative. It performs on a graph. The second one is a modified Deferred Acceptance Algorithm (DAA) using counter-proposals. It finds a many-to-one matching. We illustrate with the example of video caching from a content creator’s servers to a service provider’s servers.

\section{INTRODUCTION}
Game theorists have been interested in looking at the stability of the structures resulting from the coalition formation process and have shown the importance of the concept in real-life applications\textsuperscript{1}. The stability exhibits respect of the individual preferences and incentives. An unstable structure would result in deviations [3]. Coalition and matching games open the way through an important number of interesting applications, notably in wireless networks where much work remain to be done in the decentralized decision taking paradigm. These game-theoretic tools allow for more tractability in the formulations and a reduced complexity. Nevertheless, some of the commonly used assumptions (substitutability, responsive preferences, see [3]) limit the modeling to the cases without complementarities or peer effects. Recent theoretical works [9] (and references therein) overcome these difficulties and have paved the way through a suitable modeling of systems with complementarities and peer effects. As examples of such systems in wireless networks we have WiFi and its related anomaly, Device-2-Device with multi-hop relaying, virtual MIMO exploiting multi-users diversity, etc.

\section{RELATED WORKS}
Much work has been devoted to super-additive coalition games (where the grand-coalition forms [2]) and to the stabilization of structures by payoff distribution. Nevertheless, real-word problems may not verify the super-additivity assumption. More recently, researchers have been interested in

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\textsuperscript{2}The coalition structure graph is defined as the graph $G = (\mathcal{S}, \mathcal{E})$, where the set of vertices is the set of coalition structures and the set of edges is defined such that it exists an edge $e = (S, S')$ between vertices $S$ and $S'$ if $S'$ can be obtained from $S$ by merging two coalitions of $S$ or by splitting a coalition of $S$ into two disjoint ones. By definition, the graph is of size $\mathcal{B}(N)$ (Bell number).
The first algorithm is a centralized algorithm enumerating the set of core stable structures. It performs on a graph called the compatibility graph and does not require a tie-breaking rule (indifference between payoffs). The second algorithm is a decentralized algorithm finding a many-to-one core stable matching. It is a modified DAA that uses the pairwise alignment of the preferences over groups to reduce the preferences of the proposers over individuals.

5. POTENTIAL COALITION GAMES AND VIDEO CACHING

5.1 Potentials coalition games

Let $\Gamma = (N, v, \{u_i\}_{i \in N})$ define a coalition game in characteristic form. The set $N$ denotes the set of players of cardinality $N$, $\{u_i\}_{i \in N}$ denotes the set of their individual utilities and $v : N \to \mathbb{R}$ is the characteristic function of the game. We define the set of coalitions $C$. The potential function $\Phi$ is an ordinal coalition potential for the game $\Gamma$ if for every player $i \in N$,

$$u_i(C) > u_i(C') \text{ iff } \Phi(C) > \Phi(C'), \text{ for every } C, C' \in C \quad (1)$$

where $u_i(C)$ is player $i$’s utility of the payoff it receives when taking part in the coalition $C$.

A coalition game in characteristic form with set of players $N$ admitting an ordinal potential $\Phi$ is called an ordinal coalition potential game in characteristic form and denoted,

$$\Gamma = (N, v, \{u_i\}_{i \in N}, \Phi) \quad (2)$$

5.2 An application: the video caching

As an example consider a video caching matching market between a content creator $P$ and a service provider $S$. The content creator $P$ has a set $L$ of $L$ videos. The service provider $S$ has a set $R$ of $R$ caching servers. The caching servers may differ in the Quality of Service (QoS) for the creator’s videos and a single server. We assume that the service provider can compute the impact of caching servers. The caching servers may differ in the Quality of Service (QoS) for the creator’s videos and a single server. We assume that the service provider can compute the impact of caching over the number of views of the videos in the coalition as the following fixed point equations,

$$n_l(C) = n_l(1)\gamma_r + \sum_{l' \subseteq C \cap L} a_{ll'} n_l(C')$$

where $a_{ll'}$ is an impact factor of $l$ over $l'$ and $\gamma_r \in T$ is QoS gain obtained by caching in the server $r$. We define the characteristic function $v : C \to \mathbb{R}$ of a coalition as,

$$v(C) = \sum_{l \in C \cap L} n_l(C)\alpha$$

where $\alpha$ is the constant monetary income generated by a view of a video. We now assume $v(C)$ is shared among the players in $C$ via a Nash bargaining (null threats) with concave individual utilities $u_l(x_l) = x_l^{\alpha_p}$ for the content creator’s videos and $u_r(x_r) = x_r^{\alpha_p}$ for the service provider’s server $r$. The servers’ bargaining powers increase in their factor of quality. We obtain that the individual payoffs increase in $\chi C = v(C)$. The set of core stable structures. It performs on a graph $G$ and denoted, $\Gamma = \sum_{l \in C \cap L} n_l(C)\alpha$, which is called the set of core stable structures.

6. ANYTIME ENUMERATIVE CLIQUES-BASED ALGORITHM

In this section we introduce an anytime cliques-based algorithm for enumerating the core stable coalition structures in the case of an ordinal coalition potential game $\Gamma$ with a set of coalitions $C$. The algorithm has been constructed over the game-theoretic sequential description of convergence to a stable structure provided in [9].

Define the coalitions weighted-vertices undirected graph $(G, \Phi)$ such that the graph $G = (V, E)$ is the complement of the intersection graph of the coalitions in $C$. The set of vertices $V$ is in bijection with the set of coalitions $C$ but for the sake of simplicity and clarity we will further identify a coalition $C_i \in C$ with the vertex $v_i \in V$ of the graph. Let $V' \subset V$ be a subset of the nodes, we denote $G' = (V', E(V'))$ the subgraph induced by $V'$. Each vertex $v_i \in V$ is weighted by the value $\Phi(C_i)$ and the set of edges $E$ is defined by the adjacency matrix $A = (a_{ij})_{(i,j)\in E}$ such that,

$$a_{ij} = \begin{cases} 1 & \text{if } C_i \neq C_j \text{ and } C_i \cap C_j = \emptyset \\ 0 & \text{if } C_i = C_j \text{ or } C_i \cap C_j \neq \emptyset. \end{cases}$$

A clique $\sigma$ of the graph $G = (V, E)$ is a complete subgraph of $G$. In case the clique $\sigma$ may not be included in a superior sized clique $\sigma'$ without losing the completeness property, $\sigma$ is said to be maximal. If $\sigma$ is the clique with the highest cardinality (in terms of number of vertices $|\sigma|$) the clique is maximum. For the sake of clarity and simplicity, we will further identify a clique $\sigma$ with the set of its coalitions given by $C(\sigma)$. Any two coalitions $C_1$ and $C_2$ will be called compatible if they have no common players, i.e. $C_1 \cap C_2 = \emptyset$. In other words, if the nodes $v_1$ and $v_2$ are adjacent. More generally, the coalitions of $C' \subseteq C$ are compatible if the set of their nodes form a clique in the graph $G$. Denote $G_1$ the subgraph induced by the set of nodes of $k^{th}$ highest weight. As an example, $G_1$ is the subgraph induced by the nodes with maximum weights, i.e. the subgraph such that any node in $G_1$ is a coalition given each player its most preferred achievable share in the game.

The algorithm iteratively builds a forest (i.e. a set of disconnected trees) such that each tree is rooted by one of the maximal clique obtained at initialization and each vertex in each tree is a stable structure. A vertex that is not a leaf is a subgame stable structure (this is the anytime property of the algorithm) and a leaf is a stable structure for the game.
A DECENTRALIZED ALGORITHM: BACKWARD DEFERRED ACCEPTANCE

We now show that a modified version of the Gale and Shapley’s Deferred Acceptance Algorithm (DAA) in its college-admission form with servers’ preferences over groups of videos and videos’ preferences over individual servers is a core stable matching mechanism for the many-to-one matching games with complementarities, peer effects and pairwise alignment of the preferences (see Algorithm 2).

The Backward Deferred Acceptance Algorithm (BDAA) is similar to the DAA in many aspects. It involves two sets of players that have to be matched. Every player from one side has a set of unacceptable players from the other side. In our case, a server and a video are always acceptable. As in DAA, the algorithm proceeds by proposals and corresponding acceptances or rejections. The main difference resides in the notion of counter-proposals, introduced to tackle the problem of complementarities.

To go into more details: Knowing the mutual impact factors of the videos and their own QoS gains, the servers compute all the possible coalitions they can form and the corresponding payoff vectors. They can thus build their preference lists over groups (Steps 1b). Then, every server transmits to each of its acceptable videos the maximum achievable payoff it can achieve in the coalitions it can form with the server (Step 1c). Every video can thus build its reduced list of preferences over individual servers: if the video prefers server $r_j$ to server $r_i$ if the maximum achievable throughput with $r_j$ is strictly greater than its maximum achievable throughput with $r_i$ (Step 1d). BDAA then proceeds by rounds during which videos make proposals, servers make counter-proposals and videos accept or reject (from Step 2a to Step 2h). $L(r)$ is the list of all videos that have proposed at least once to server $r$. $L^*(r)$ is a dynamic list that is reinitialized to $L(r)$ at the beginning of every new round (Step 2b). In each round of the algorithm, every unengaged video proposes to its most preferred server for which it has not yet proposed (Step 2a). Every server receiving proposals adds the proposing players to its cumulated list of proposers and reinitializes its dynamic list (Step 2b). Using $P^*(r)$ it then searches for its most preferred coalition involving only videos from the dynamic list and emits a counter-proposal to these videos. This counter-proposal contains the payoff the videos can achieve in this coalition (Step 2c). Each video compares the counter-proposals it just received with the best achievable payoffs obtained with the servers it has not proposed to yet (Step 2d). If one of these best achievable payoffs is strictly greater than the best counter-proposal, the video rejects the counter-proposals and continues proposing (Step 2d, Step 2h). Otherwise, the video accepts its most preferred counter-proposal (Step 2d). Given a counter-proposal, if all the videos accept it, then they are engaged to the server (Step 2e). If at least one video does not agree, then the
Algorithm 2: Backward Deferred Acceptance

Data: For each server: The set of acceptable videos. For each video: The set of acceptable servers.
Result: A core stable structure S

1 begin
2 Step 1: Initialization;
3 Step 1.a: All servers and videos are marked unengaged. L(f) = L∗(f) = ∅, ∀f;
4 Step 1.b: Every server r computes possible coalitions with its acceptable videos, the respective videos’ payoffs and emits its preference list P#(r);
5 Step 1.c: Every server r transmits to its acceptable users the highest payoff they can achieve in coalitions involving r;
6 Step 1.d: Every video l emits its reduced list of preference P′(l);
7 Step 2 (BDAA);
8 Step 2.a, Videos proposals: According to P′(v), every unengaged video v proposes to its most preferred acceptable server for which it has not yet proposed;
9 Step 2.b, Lists update: Every server r updates its list with the set of its proposers:
10 L(r) ←− L(r) ∪ {proposers} and L∗(r) ←− L(r);
11 Step 2.c, Counter-proposals: Every server r computes the set of coalitions with users in the dynamic list L∗(r) and counter-proposes to the videos of their most preferred coalition according to P#(r);
12 Step 2.d, Acceptance/Rejections: Based on these counter-proposals and the best achievable payoffs offered by the servers in Step 1.c to which they have not yet proposed, videos accept or reject the counter-proposals;
13 Step 2.e: If all videos of the most preferred coalition accept the counter-proposal of an server r, all these videos and r are marked engaged;
14 Step 2.f: Every unengaged server r updates its dynamic list by removing videos both having rejected the counter-proposal and being engaged to another server: L∗(r) ←− L∗(r) \ {engaged rejecters};
15 Step 2.g: Go to Step 2.c while there are unengaged videos with L∗(r) ≠ ∅;
16 Step 2.h: Go to Step 2.a while there are unengaged videos that can propose.

server is unengaged (Step 2.e), it updates its dynamic list by removing the videos both having rejected its counter-proposal (Step 2.f) and being engaged with an other server. The counter-proposals continues up to the point when no unengaged server can emit any counter-proposal (Step 2.g). The current round ends and the algorithm enters a new round (Step 2.h). The algorithm stops when no more videos are rejected (Step 2.h). A stable matching is obtained.

Proposition 4 In a many-to-one matching game with complementarities, peer effects and pairwise alignment of the preferences, the Backward Deferred Acceptance Algorithm (BDAA) is a core stable matching mechanism.

Assume R servers and L videos. At each round, the number of counter-propositions by servers is at most R. From the pairwise alignment of the preferences, at each step there is at least one server whose counter-proposal has been accepted. Since the algorithm stops if no more videos are rejected or all rejected videos have already proposed to all their available servers, then in at most \(L^2 - 2L + 1\), every non-engaged mobiles have proposed to all servers. Hence the complexity of our algorithm is \(O(n^2)\) where \(n = \max(R, L)\).

8. CONCLUSION

We have shown two new algorithms for finding core stable structures. The first algorithm is a centralized anytime enumerative one and the second is a decentralized deferred acceptance based one. We need to go further in the analysis of the complexity of the enumerative algorithm and search for the requirements over the coalitions compatibility graph to make it fall in classes of graphs with low complexity in the maximal cliques enumeration.

9. REFERENCES